

AI FOR HACKER

Automatic Exploit Generation for Application Source Code Analysis

THE TEAM

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TO ANALYZE ~ 400 APPLICATIONS...



WE NEED THE AI



THERE ARE DIFFERENT KINDS OF ROBOTS

Marketing approach

- Interactive Application Security Testing (IAST)
- Dynamic Application Security Testing (DAST)
- Static Application Security Testing (SAST)
- Technical approach
- Black Box/White Box
- Static/Dynamic
- Scientific approach It's all relative



DAST

We don't have access to [server] application Fuzzing/Fault injection

Pro

- Easy to implement/Easy to verify results/Low level of false positives
- Language/Framework/Backend independent

Cons

- Weak API coverage/Auth/Web 2.0
- Application should be deployed/Can terminate app*
- $(O(c^n), c > 1)$ **

*And admins will terminate you **Never stops





SAST

We have access full access to application [source code]

Model checking/correctness properties of finite-state systems

Pro

- [possible] Good coverage/Don't need to deploy app
- [possible] Good performance*

Cons

- Hard to implement/Hard to verify results
- [can generate]a lot of false positives/Language dependent
- K := { (i, x) | program i will eventually halt if run with input x} *

*Because of computation timeouts

**The halting problem



SAST

Report View 🔀									
Γ	🗆 🗭 Fi	ndings (57 417)	1	57	- 64	171	Vulnerability Type		
L	÷ 🧕	A1 - Injection (3 725)	ж.	a_{L}	- 12	4 C J -	rossSiteScripting		
		A2 - Cross Site Scripting (XSS) (1 616)	- 21				rossSiteScripting		
L	H 😣	A3 - Broken Authentication and Session Managem	0	-	r ngi r	17001	crossSiteScripting		
L	8	A4 - Insecure Direct Object Reference (255)		22	High	Type I	CrossSiteScripting		
L	2	A5 - Cross Site Request Forgery (CSRF) (38)			High	Type I	CrossSiteScripting		
L	4	A6 - Security Misconfiguration (18 282)		20	High	Type I	CrossSiteScripting		
L	H 💈	A7 - Insecure Cryptographic Storage (39)	-	22	High	Type I	CrossSiteScripting		
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1					1 · · · # · ·	** ** * *			
١.,	803	if(\$this->do_debug >	if(\$this->do_debug >= 1) {						
Ιđ	804	echo "SHTP -> ERROR: " . \$this->error["error"] .							
	805	": " .	ly . \$t	y . \$this->CRLF;					
	806	}							
	807	return false;							
	808	}							

SAST



IAST

Have full access to application [source code]/system and can patch it

Fuzzing/Instrumentation/Data [control] flow tracing

Pro

- Can combine strengths of SAST and DAST
- Control of dataflow/Second chance vulns/binary analysi

Cons

- Can combine weaknesses of SAST and DAST
- Need fuzzer/Need to patch server
- Generates tons of results (execution trace)
- Need to have/patch "live" system





CAN WE USE (.AST)

Dynamic Application Security Testing (DAST) and Static Application Security Testing (SAST) -- Issue Type Coverage





URL-TO-SOURCE MAPPINGS

SAST and DAST have produces incompatible output SAST: line of code, CFG DAST: Input data (HTTP Request)



HYBRID ANALYSIS!

Real-Time Hybrid Correlation





http://bit.ly/17wbvnL



REALITY

Need to have and to patch "live" system/source code

Need to analyze application several times

Magic to correlate "line number" (SAST) and "input data" (DAST) $(O(c^n), c > 1)$



PERFECTION?

No live system Low level of false positives Automatic exploits generation!





PERFECTION: NO LIVE SYSTEM

Need to use static analysis

Proper model representation is half the battle



ABSTRACT SYNTAX TREE while return condition variable compare body op: ≠ name: a while b != 0 if a > bvariable constant branch name: b value: 0 a := a - b else condition lif-body else-body b := b - a compare assign assign return a op: >variable variable variable variable bin op bin op name: a name: a op: name: b op: name: b variable variable variable variable name: a name: a name: b name: b

CONTROL FLOW GRAPH



PROGRAM DEPENDENCE GRAPH





SYMBOLIC EXECUTION



SYMBOLIC EXECUTION!

Microsoft Automata

KLEE/Kleaver



SYMBOLIC EXECUTION :(

Path Explosion *

Full support of language (functions/frameworks/environment)** [sometimes] too far from real code [execution flow]***

*Number of paths grows exponentially with program size and can be infinite ****

Zillions man-hours with endless updates**

***SAT was the first known NP-complete problem, as proved by Stephen Cook in

**** Never stops



!FALSE POSITIVES == EXPLOITS



EXPLOIT IS USEFUL TO

prove that vulnerability exists* make additional [dynamic | automatic] checks** create test cases for QA generate signatures/virtual patches for AF/IDS***

* get devs to shut up and fix the bug
**automatic verification via fuzzing
***self-defending application







Languages grammar

- Input functions
- Filtering functions

Potentially Vulnerable Functions (PVF)

- Related Vulnerabilities
- Related Preconditions (Bad Inputs)
- Related Exploit Creation Rules

Safe functions

Can be called without any risk

WHY SLICING?

Problem Occurred	Reportir Lines Before Lines After:
'UpdateValues' has encountered a problem. null OK Details >>	▼ Notes
SmartHeap Library SmartHeap Library	
ОК Отмена	



```
# Hint 1.
$inc = './inc/';
```

```
$file = 'config-for-this-site';
```

```
include $inc.$file.".php";
```

```
# Hint 2.
echo stripslashes("<script>alert(1)</script>");
```



<?php // /test.php

SOLVER

print base64_decode (\$_GET['x']) ;

?>

exploit:

GET

/test.php?x=PHNjcmlwdD5hbGVydCgxKTwvc2NyaXB0Pg%3D %3D



DEMO

INSIDE IN [ISLAND] GRAMMAR

SELECT CONCAT (last_name, ', ', first_name) full_name FROM mytable ORDER BY full_name
SELECT t1.name, t2.salary FROM employee AS t1, info AS t2 WHERE t1.name = t2.name;
SELECT t1.name, t2.salary FROM employee t1, info t2 WHERE t1.name = t2.name;
SELECT college, region, seed FROM tournament ORDER BY region, seed;
SELECT college, region AS r, seed AS s FROM tournament ORDER BY r, s;
SELECT college, region, seed FROM tournament ORDER BY 2, 3;
SELECT t1.name, t1.name, t2.salary FROM employee WHERE id = \$i;
SELECT * FROM foo ORDER BY RAND (NOW ()) LIMIT 1;

Change MySQL Grammar

"SELECT t1.name, t1.name, t2.salary FROM employee WHERE id = \$i"
\$i=1+union+select+1,2,3--+ // SQLi Exploit!

CONDITIONS

We can't [symbolically | interactive] resolve all part of equation

Session id's in files:

(file('../admin/conf/config.inc')[2] == session_id())

Session values are set:

\$_SESSION["admin_login"]==true

External connections:

ftp_connect(str_replace('ftp://', ", \$_POST['ftpsite']))

Configuration:

- !((strpos(php_sapi_name(), 'apache') !== False))
- sqlsrv_connect('***', array('Database' => '', 'UID' => '***', 'PWD' => '***'))==True

BACKDOORS?

Exploit:

GET/core/jscss.php?files=%2F..%2F..%2Fetc%2Fpasswd

Conditions:

(md5(\$_GET['PA']) === 'bb2a4974d7aca7da8735c70371361c0f')



BACKDOORS!

...we use it for emergency support cases when we need to access files but we don't

have a password....HOBOCTW LEGO Edition SP



🦆 ЦИТАТА

DEMO

PRACTICAL TESTS

SECOND CHANCE?

Cross Site Scripting Vulnerability

Exploit: GET /viewResults.php HTTP/1.1 Code: print \$question . "
";



Condition

(mysql_fetch_assoc(mysql_query(('SELECT * FROM tblquestions, answers WHERE tblquestions.QID = answers.QID AND answers.QID = \'' . \$_GET['h1'] . '\'')))['Question'] === '<script>alert(1)</script>')

SECOND CHANCE!





CONCLUSIONS

Exploit generation can improve .AST

- Reduce false positive
- Add transparency
- Helps o hack stuff

Condition resolver can help do detect

- Authentication condition and access control issues
- Hidden execution paths (e.g. backdoors)
- Hardcoded conditions

Combination of symbolic and real execution is useful

- Reduce labor input
- Improve performance
- Helps to balance CPU/time/memory



RELATED WORKS

Chandrasekhar Boyapati, Paul Darga. Eficient software model checking of data structure properties.

Keshav Pingali, Micah Beck, Richard Johnson, Mayan Moudgill, and Paul Stodghill. Dependence flow graphs: an algebraic approach to program dependencies.

E. Morel and C. Renvoise. Global optimization by suppression of partial redundancies.

R. Cytron, J. Ferrante, B. K. Rosen, M. N. Wegman, and F. K. Zadeck. An eficient method of computing static single assignment form.

Vugranam C. Sreedhar and Guang R. Gao. Computing u-nodes in linear time using dj-graphs

Mark N. Wegman and F. Kenneth Zadeck. Constant propagation with conditional branches.

Ron K. Cytron and Jeanne Ferrante. Eficiently computing u-nodes on-the-fly.

Thomas Ball, Rupak Majumdar, Todd D. Millstein, and Sriram K. Rajamani. Automatic predicate abstraction of c programs.

Thomas Ball and Sriram K. Rajamani. Bebop: A symbolic model checker for boolean programs.

David Binkley. Interprocedural constant propagation using dependence graphs and a data-flow model.

AEG

MAYHEM

The Essence of Command Injection Attacks in Web Applications, http://www.cs.ucdavis.edu/~su/publications/popl06.pdf

http://qspace.library.queensu.ca/bitstream/1974/5651/3/Alalfi_Manar_H_2010April_PhD.pdf

SPECIAL THANKS

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